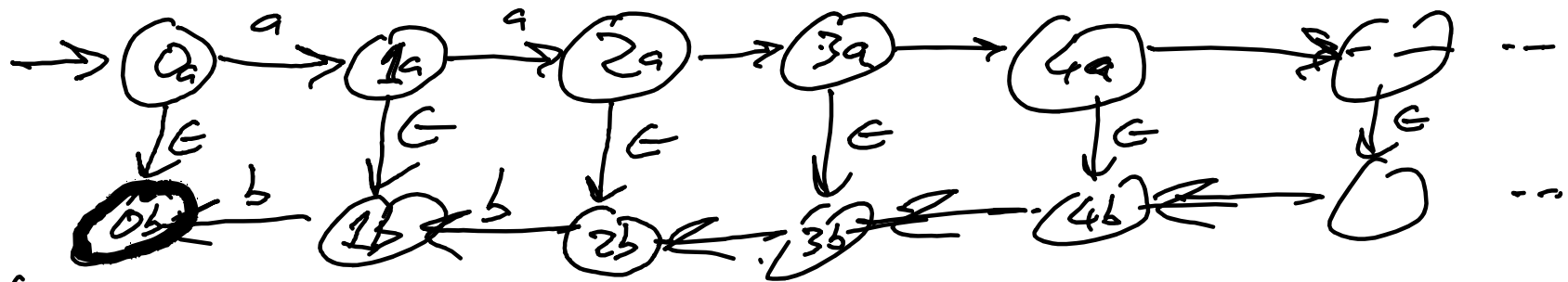


$$L = \{a^n b^n : n \geq 0\}.$$

Lecture 9

Fictional infinite state NFA:



(Motivates PDA in lecture slides)

Context-free grammar:

$$S \rightarrow \epsilon \mid a S b$$

Random generation:



Why could we not write?

$$S \rightarrow \epsilon \mid a A b B$$

$$A \rightarrow \epsilon \mid a A$$

$$B \rightarrow \epsilon \mid b B$$

No "context" links them

$$L = \{a^n b^m : n \geq m \geq 0\}$$

$S \rightarrow \text{"Some As"} \text{"Matching"}$
 $\text{"Some As"} \rightarrow \epsilon \mid a \text{"Some As"}$
 $\text{"Matching"} \rightarrow \epsilon \mid a \text{"Matching"} \mid b$

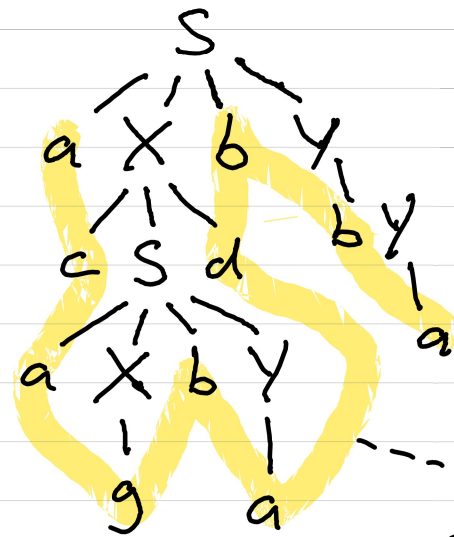
Example context-free grammar

$$S \rightarrow aXbY$$
$$X \rightarrow cSd$$
$$X \rightarrow g$$
$$Y \rightarrow bY$$
$$Y \rightarrow a$$

Example word generation

$S \rightarrow aXbY$
 $\rightarrow aXbbY$
 $\rightarrow acSdbbY$
 $\rightarrow acSdbba$
 $\rightarrow acacXbYdbba$
 $\rightarrow acagbYdbba$
 $\rightarrow acagbadbba$

Derivation tree



--- Non-terminal symbols found during in-order traversal of derivation tree